

# Expressions

- operands & operators
- precedence rules
  - Java: nested first, then precedence, then L→R

Java operator precedence (highest at top)

```

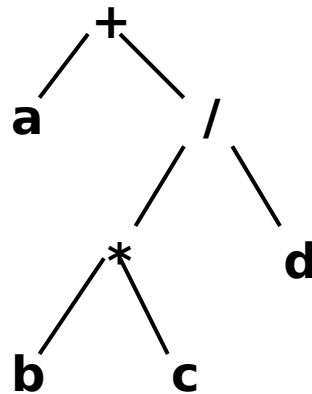
!
*, /, %
+, -
>, <, >=, <=
==, !=
&&
||
=

```

- APL: all operators have the same precedence, R→L
- Pascal: fewer levels of precedence and Boolean operators have higher precedence than the relational operators:

```
if (a=b) | c then ... ;
```

- expression tree, e.g.  $a+b*c/d$



- side effects
  - function calls
- operator vs function
- operator forms
  - binary (dyadic), unary (monadic), trinary (tridaic)
  - overloading

# Boolean Expressions

- logical values
- relational operators
- Boolean operators
  - and, or, not
- precedence
  - Ada vs Pascal

`a > b and c + 4 < d`

- short-circuit

`if (i < 0) or (a[i] > a[i+1]) then ...`

- operators

Java	Ada
<code>&amp;&amp;</code>	<code>andthen</code>
<code>  </code>	<code>orelse</code>
<code>&amp;</code>	<code>and</code>
<code> </code>	<code>or</code>

# Mixed-mode Expressions

- operands of different types

```
var a : integer;
```

```
var b : float;
```

```
:
```

```
... a + b ...
```


- strongly-typed → not legal, e.g. Ada
  - firmly-typed → widening conversion occurs, e.g. Java
  - weakly-typed → most conversions occur, e.g. C
- casting

```
... Float(a) + b ...
```

# Statements

- change state and perform actions
- assignment statement
- conditional statement
- looping statement
- procedure call statement
- Algol 68 - no difference between statement and expression

# Assignment Statement

- l-value vs r-value
- restrictions on l-value
  - variable vs expression
- conversions
  - strongly-typed vs firmly-typed vs weakly-typed
  - widening vs narrowing
  - casting
- assignment as operator
  - Compound assignment in C, C++ and Java (`+=`, `-=`, `*=`, `/=`)  
`a += expression;      VS      a = a + expression;`
  - increment and decrement operator (`++`, `--`)  
`a = 1; b = ++a;      VS      a = 1; b = a++;`  


- multiple assignment
  - assignment as expression (Algol 68, C, C++, Java)  
`a = b = c = expression;`
  - lists or pairs as a left-hand side (Perl)  
`($a, $b) = ($b, $a);`

# Compound Statement

- where control structures allow only a statement
  - e.g. Pascal, Java

```
if (a < b)      and
    c = d;
```

```
if (a < b) {
    c = d
    e = f;
}
```

- explicit statement terminators
  - e.g. Ada

```
if (a < b) then c := d; e := f; end if;
```

- advantages
- vs blocks
  - declarations

# Selection Statements - If

- FORTRAN
  - arithmetic if `IF (C1) L1, L2, L3`
  - logical if `IF (logical expression) statement`
  - implied goto
- ALGOL 60
  - single statement
  - compound statement
  - single entry, single exit - structures statements
  - dangling else (Pascal, C, Java)  
`if C1 then if C2 then S1 else S2`
- Ada if
  - statement terminator
  - `elsif`

# Conditional Operator

- There's a ternary operator offered by some languages that's commonly called a 'conditional operator' (or an 'inline-if')
  - I don't feel like making a slide about it. So remind me to do it on the board, k?

# Selection Statements - Case

- early versions - computed GOTO (FORTRAN) `GOTO (10,20,30) I`
- first implemented in ALGOL W
- Pascal
  - case month of
    - discrete type `1, 3, 5, 7, 8, 10, 12 : days := 31;`
    - `4, 6, 9, 11 : days := 30;`
    - `2 : if years mod 4 = 0`
    - `then days := 29`
    - `else days := 28`
    - `end`
- Ada
  - case ch is
    - ranges of values `when '0' .. '9' => put_line("digit");`
    - default condition `when 'A' .. 'Z' => put_line("letter");`
    - `when others => put_line("special");`
    - `end case;`

- C, C++, Java
  - switch statement
  - closer to computed goto
  - use of break

```
switch (month) {
    case 4:
    case 6:
    case 9:
    case 11: days = 30;
             break;
    case 2:  if (years % 4 == 0) days = 29; else days =
             28;
             break;
    default: days = 31;
             break;
}
```

# Conditional Loop

- while statement
  - pre-test loop
  - compound statement vs terminator
  - `break` or `exit`
- repeat statement
  - post-test loop
  - Pascal vs C, Java
- loop statement
  - Ada - infinite loop
  - `exit` or `break`
  - in-test loop
  - multi-exit loop

```
sum := 0;
    loop
        get(number);
        exit when number < 0;
        sum := sum + number;
    end loop;
```

# Fixed Iteration Loop

- DO statement (FORTRAN)
- for statement
  - fixed increment - Pascal, Ada

```
for cv in low .. high loop ... end loop;
```
  - generalized increment - ALGOL 60 (index can be real)

```
for i := 0.0 step 0.1 until 6.0 do S
```
- final index value
  - first that fails vs last that passes vs undefined
  - scope of index

```
for (int i = 0; i < n; i++) S
```

- predetermination of number of repetitions
  - test and final value expressions
  - changes to index within loop
  - premature exit from loop
  - e.g.

```
j := 2;  
for i := 1 to 10*j step 7-j do  
    j := j + 1;  
end;
```

- Pascal (final value and increment are computed once)

i	j	final value	increment	new j
1	2	20	5	3
6	3	20	5	4
11	4	20	5	5
16	5	20	5	6

- Algol60, C, C++, Java (final value and increment are computed in every iteration, negative increment is decrement)

i	j	final value	increment	new j
1	2	20	5	3
6	3	30	4	4
10	4	40	3	5
13	5	50	2	6
15	6	60	1	7
16	7	70	0	8
16	8	80	-1	9
15	9	90	-2	10

- C, C++ & Java for loop

- really a while loop

```
for (e1; e2 ; e3) S  ⇔      e1;  
                        while (e2) {S; e3}
```

- may lead to unreadable programs

```
for(i = f = 1; i < n; f *= ++i) {}
```

- what is computed by this program?

# Iterator

- iterate over all elements of a data structure
- e.g. Java 1.5

```
for ( VarDcl : Expression ) Statement
```

- where `Expression` is an array or an object implementing `Iterable`
- the statement is executed with the variable referencing the elements of the array or data structure in no specified order
- e.g.

```
int a[];      List l;  
a = new int[100];  l = new List();  
      :           :  
for ( int e : a ) { for ( Object o : l ) {  
    sum = sum + e;      out.println(o);  
};      };
```

## Goto Statement

- heavily used in early languages
  - machine language branch instruction
- Dijkstra's "Goto Statement Considered Harmful"
  - reasoning about programs
  - structured programming
- restricted `goto`'s
  - `break`, `exit`
    - with label
  - `continue`
  - exception handling

Example Time!

# Exception Handling

- exception
  - exceptional but not necessarily unexpected event
  - unexpected event
- handling exceptions
  - up to programmer, else crash
  - language support
- exception handling
  - PL/I, Ada, Java, C++
  - event occurs, exception raised (thrown), current unit terminated, exception handler invoked
- user defined exceptions

- in Ada & Java
  - handler attached to block (frame, try statement)
  - replaces execution of block
  - if not caught, is propagated to caller

## Ada

---

```
begin
  ...
  if a(i) < 0 then ...
  ...
exception
  when constraint_error => ...
end;
```

## Java

```
try {
  ...
  if (a[i] < 0) ...
  ...
}
catch (ArrayIndexOutOfBoundsException e) {
  ... }
```

- an example (Ada)

```
loop
  begin
    get(x);
    exit;
  exception
    when data_error =>
      skip_line;
      put_line("Error in number, try again");
  end;
end loop;
```

- Java
  - checked vs unchecked exceptions
  - exceptions are objects
- C++
  - more complex and less general
- Eiffel
  - two choices
    - report error and terminate
    - fix up and retry failed operation (retry command)

- exceptions shouldn't be overused
  - only for exceptional conditions

```
type Days is
    (Sunday, Monday, Tuesday, Wednesday, Thursday, Friday, Saturday);

if today = Saturday then
    tomorrow := Sunday;
else
    tomorrow := Days'succ(today);
begin
    tomorrow := Days'succ(today);
exception
    when constraint_error =>
        tomorrow := Sunday;
end;
```

# Assertions?

- If I look bored, ask me about assertions

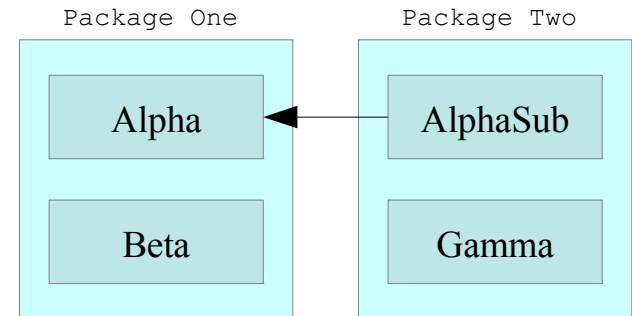
# Program Structure

- effects of design process on language design
- procedural
  - stepwise refinement or top-down design
  - control structures
  - substeps are procedures
- abstraction
  - hide unwanted detail to control complexity
  - procedural abstraction
    - hide implementation
- encapsulation
  - prevent or control access to what is hidden
  - interface defines what is visible
    - signature for procedure

# (Encapsulation)

- When it comes to Information Hiding, some languages (e.g. Java) permit a decent level of control
- Consider the following example from Oracle's Java tutorial page ( <http://download.oracle.com/javase/tutorial/java/java00/accesscontrol.html> )

Visibility of Alpha's Members				
Modifier	Alpha	Beta	AlphaSub	Gamma
public	Yes	Yes	Yes	Yes
protected	Yes	Yes	Yes	No
(none)	Yes	Yes	No	No
private	Yes	No	No	No



- procedural abstraction doesn't scale up
  - teams of programmers
- modules
  - collections of related procedures and methods
  - encapsulation with well-defined interface
  - raises the level of abstraction
  - in object-oriented languages class is a module
- classes
  - represent real-world entities
  - same entities in analysis, design and implementation
  - attributes, services and behavior
  - hidden state, public interface
  - multiple instances
  - object identity
  - interaction
    - client object and server object

# Abstract Data Types

- built-in types
  - set of values and set of operations
  - don't need to know representation (encapsulation)
- set of types defines applicability of language
  - add more built-in types
  - define new types - extension
    - declare type and operations
    - hide representation and implementation
    - abstract data type
- module vs class
  - contains a type vs defines a type

# Modifiability and Reuse

- modifiability
  - large systems
  - limit scope of change
  - loosely-coupled modules
- reuse
  - unit of reuse
    - procedure
    - module
    - class

# Procedural vs Object-Oriented

- procedure as separate entity or part of class or both
- function vs procedure
  - change state vs product a value
  - pure function
  - don't differentiate - Algol 68 & C++
- methods
  - aka procedures, member functions

# Java

```
class BankAccount {
    // private attribute declarations
    ...
    public BankAccount() {
        ...
    } // constructor
    public void deposit(int amount) {
        ...
        } //deposit
    public int getBalance() {
        ...
    } // getBalance
} // BankAccount
```

# Java

- declaration and creation

```
BankAccount bk1 = new BankAccount();  
BankAccount bk2 = new BankAccount();
```

- method calls

```
bk1.deposit(6);  
int am1 = bk1.getBalance();
```

# Pascal

```
program BankAccountEx(input, output);
  type BankAccount = ...
  var bk1, bk2 : BankAccount;
      am1 : Integer;
  ...
  procedure makeBankAccount(var b : BankAccount);
  begin ... end {makeBankAccount};

  procedure deposit(var b : BankAccount; amount : Integer);
  begin ... end {deposit};

  function getBalance(b : BankAccount) : Integer;
  begin ... end {getBalance};

begin
  makeBankAccount(bk1);
  makeBankAccount(bk2);
  ... deposit(bk1, 6);
  ... am1 := getBalance(bk1); ...
end.
```

# Pascal

- monolithic
- type declarations
- initialization procedure
- “self” parameter
- no connection between type and procedures for type
- no hiding of representation of type

# Ada

```
package BankAccounts is
  type BankAccount is private;
  procedure makeBankAccount(b : out BankAccount);
  procedure deposit(b : in out BankAccount;
                   amount : in Integer);
  function getBalance(b : BankAccount) return Integer;
private
  type BankAccount is ...
end BankAccounts;
```

```
package body BankAccounts is
  -- definitions of makeBankAccount, deposit and getBalance
end BankAccounts;
```

# Ada

- exports `BankAccount` type and procedures
- specification defines type and procedure signatures
- body gives implementation
- representation given in private part of specification