

Topological Sort Algorithm

- Given a graph G with n vertices, generate a list L of vertices arranged in topological order.

```
// set indegree for each vertex
for(i = 0; i < n; i++)
    for(j = 0; j < n; j++)
        v[j].indegree += A[i][j];
// apply topological sort
for(i = 0; i < n; i++)
{
    find an unvisited vertex v with no predecessors;
    if(v != null)          // there is one with no predecessors
    {
        insert v at back of L;
        v.wasVisited = true;
        decrement indegree of each vertex adjacent to v;
    }
}
```

Algorithm to find an unvisited vertex with no predecessors:

```
for(i = 0; i < n; i++)
    if((v[i].indegree == 0) && (v[i].wasVisited == false))
        return v[i];
return null;          // every vertex has a predecessor
```

Bridges of Königsberg Problem

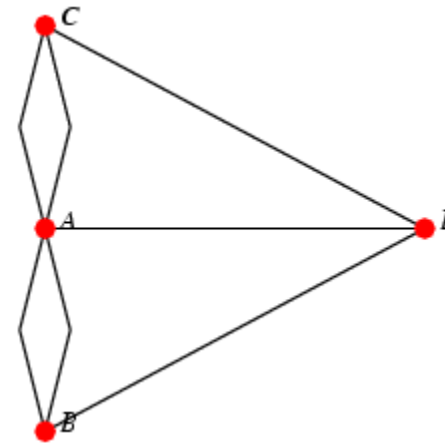
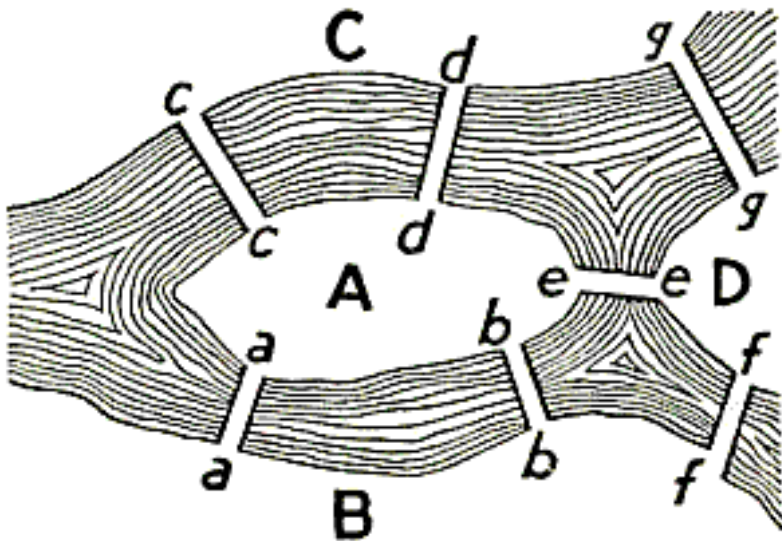
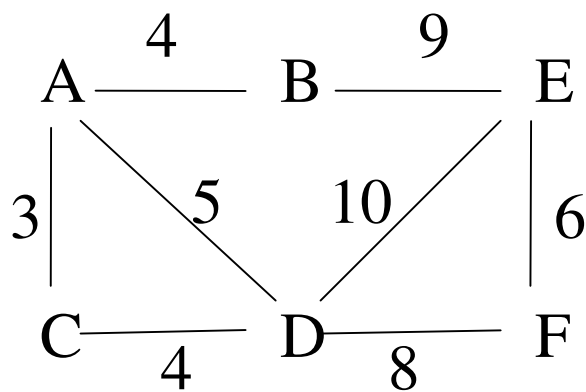


FIGURE 98. *Geographic Map:
The Königsberg Bridges.*

- There were 7 bridges in Königsberg, arranged as above (pictures from <http://mathworld.wolfram.com>).
- Is it possible to travel across every bridge exactly once, and return to the starting point?

Weighted Graphs

- Adjacency matrix representation:
 - $A[i][j] = w$ if there is an edge between vertex i and vertex j with weight w
 - $A[i][j] = \infty$ otherwise



	A	B	C	D	E	F
A	-	4	3	5	∞	∞
B	4	-	∞	∞	9	∞
C	3	∞	-	4	∞	∞
D	5	∞	4	-	10	8
E	∞	9	∞	10	-	6
F	∞	∞	∞	8	6	-

Unweighted Shortest Path Algorithm

- We wish to find the shortest path from s to all other vertices.
- Assume that for every vertex v , we have initialized $v.\text{known} = \text{false}$ and $v.\text{dist} = \infty$.

```
Q.enqueue(s);
s.dist = 0;
while Q is not empty
{
    v = Q.dequeue();
    v.known = true;
    for (each vertex w adjacent to v)
        if(w.dist > v.dist+1)
        {
            w.dist = v.dist + 1;
            w.path = v;
            Q.enqueue(w);
        }
}
```

Weighted Shortest Path Algorithm (Dijkstra's Algorithm)

- We will find the weighted shortest path from s to all other vertices.
- Assume that for every vertex v , we have initialized $v.\text{known} = \text{false}$ and $v.\text{dist} = \text{infinity}$.

```
s.dist = 0;
PQ.enqueue(s, 0);
while PQ is not empty
{
    v = PQ.dequeue();
    if(!v.known)
    {
        v.known = true;
        for (each vertex w adjacent to v)
            if(w.dist > v.dist + A[v][w])
            {
                w.dist = v.dist + A[v][w];
                w.path = v;
                PQ.enqueue(w, w.dist);
            }
    }
}
```